#### Lecture 25: Data Flow and Dependence Graphs

Kenneth M. Anderson Foundations of Software Engineering CSCI 5828 - Spring Semester, 1999

## Today's Lecture

- White-Box Testing
  - Data Flow Graphs
- Minimum Retesting

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- Program Dependence Graphs
  - Control Dependence Graphs
  - Data Dependence Graphs

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## Flow Graphs

Graph representation of control flow and data flow relationships

• Control Flow

The partial order of statement execution, as defined by the semantics of the language

• Data Flow

The flow of values from definitions of a variable to its uses

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## A Sample Ada Program to Test

function P return INTEGER is
begin
X, Y: INTEGER;
READ(X); READ(Y);
while $(X > 10)$ loop
X := X - 10;
exit when $X = 10$ ;
end loop;
if $(Y < 20)$ and then X mod 2 = 0) then
Y := Y + 20;
else
Y := Y - 20;
end if:
return $2 * X + Y;$
end P;

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## Minimizing Retesting

- Test Only What Is Affected by a Change
- Key: Dependency Analysis Also used for optimization, parallelization, ...
- At Coarse Level, Module Relationships Uses, calls, imports, includes, ...
- At Fine Level, Control and Data Flow Program dependence graphs

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## Program Dependence Graph (PDG)

- Summary Representation of "Dependence"
- Nodes Are Either Statements or Predicates or the Special Node "Entry"
- Two Kinds of Edges
  - Control dependence edge
  - Data dependence edge
- Two Subgraphs Induced by the Edges

## Control Dependence Graph (CDG)

- Informal Definition
  - For nodes X and Y in a CFG, Y is control dependent on X if, during execution, X can directly affect whether Y is executed

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## Control Dependence Graph (CDG)

- Formal Definition
  - Let X and Y be nodes in a CFG. If Y appears on every path from X to the exit node, where Y != X, then Y post-dominates X.
  - There is a control dependence from X to Y with label L iff:
    - there is a non-null path p from X to Y, starting with edge L, such that Y post-dominates every node strictly between X and Y on p; and
    - Y does not post-dominate X.

## P's Control Dependence Graph



#### P's Control Dependence Graph



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## P's Control Dependence Graph



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## Data Dependence Graph (DDG)

- Informal Definition
  - Two statements are data dependent if they might reference the same memory location and one of the references is an assignment to the memory location

## Data Dependence Graph (DDG)

- Informal Definition
  - Two statements are data dependent if they might reference the same memory location and one of the references is an assignment to the memory location
  - Intuition: If the statements cannot be switched without affecting the program, then they are data dependent

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## Data Dependence Graph (DDG)

- Formal Definition
  - Let X and Y be nodes in a CFG. There is a data dependence from X to Y with respect to a variable v iff there is a non-null path p from X to Y with no intervening definition of v and either:
    - X contains a definition of v and Y a use of v;
    - X contains a use of v and Y a definition of v; or
    - X contains a definition of v and Y a definition of v.

# P's Data Dependence Graph for X



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#### P's PDG (DDG for X Only)



#### P's PDG (DDG for X Only)



## P's PDG (DDG for X Only)



## Minimum Regression Testing

Given program P, its modified version P', and test set T used to test P, find a way, making use of T, to test P'

- Identify changes to P resulting in  $P^\prime$
- Select T', a subset of T, related to changes
  Run T' on P'

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## Goals

• Safety

Every relevant test from T must be selected

• Precision

Select only tests that exhibit different behavior

• Efficiency

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Cheap to calculate and run T'

Modifications

- Adding Statements
- Deleting Statements
- Changing Statements
- Theorem

Need only tests in T that can traverse different *regions* of statements in P and P', where regions are dependent-equivalent sub-CDGs

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## Test Selection Algorithm

procedure SelectTests Construct CDGs of P and P', with entry nodes E1, E2 T' = Compare (E1, E2)

procedure Compare (N1,N2) mark N1 and N2 visited if (children of N1 and N2 differ) then return all tests that traverse N1 else T' = NULL

for each region or predicate child C1 of N not yet visited do find C2, the corresponding child of N2 T' = T' union Compare (C1,C2) Lecture 25

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