Lecture 18: Petri-Nets (Continued)

Kenneth M. Anderson Foundations of Software Engineering CSCI 5828 - Spring Semester, 1999

Today's Lecture

- Finish the Filling Station Example
- Look at analysis techniques using Petri Nets
- Look at extensions to the basic Petri Net formalism
 - add "data" to tokens
 - add "conditionals" to transitions

Filling Station Example

- Lets model the following situation
 - Fuel Pumps
 - Spaces next to Pumps
 - A cashier that takes payment
- Questions
 - What is the concurrency that we want modeled?
 - How do we handle the parameterization of the Petri net? (e.g. lets say I want to add a pump)

Concurrency Problems

• Starvation

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Enabled transition never fired

• Deadlock

Unintended lack of enabled transitions

• V&V Tries to Detect These Problems Static and dynamic analysis techniques

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Shortcoming of Basic Petri Nets

Simplicity of building blocks leads to complexity in nets

Example: Semaphore for *n* processes requires 2n transitions and 3n+1 places

Would Like...

- *Enable* and *fire* as computations
- Tokens as data, not just control

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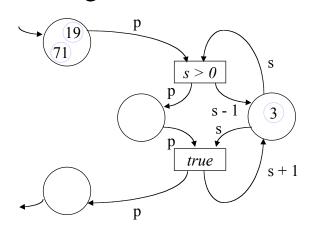
Higher-Level Petri Nets

- Some Enhancements to Basic Petri Nets
 - Typed places and information-bearing tokens
 - Predicate transitions
 - Hierarchical decomposition of places and transitions

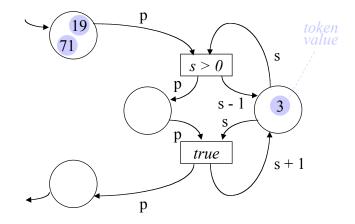
Requirement for analysis of higher-level nets: reducible to basic nets for analysis

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Higher-Level Net



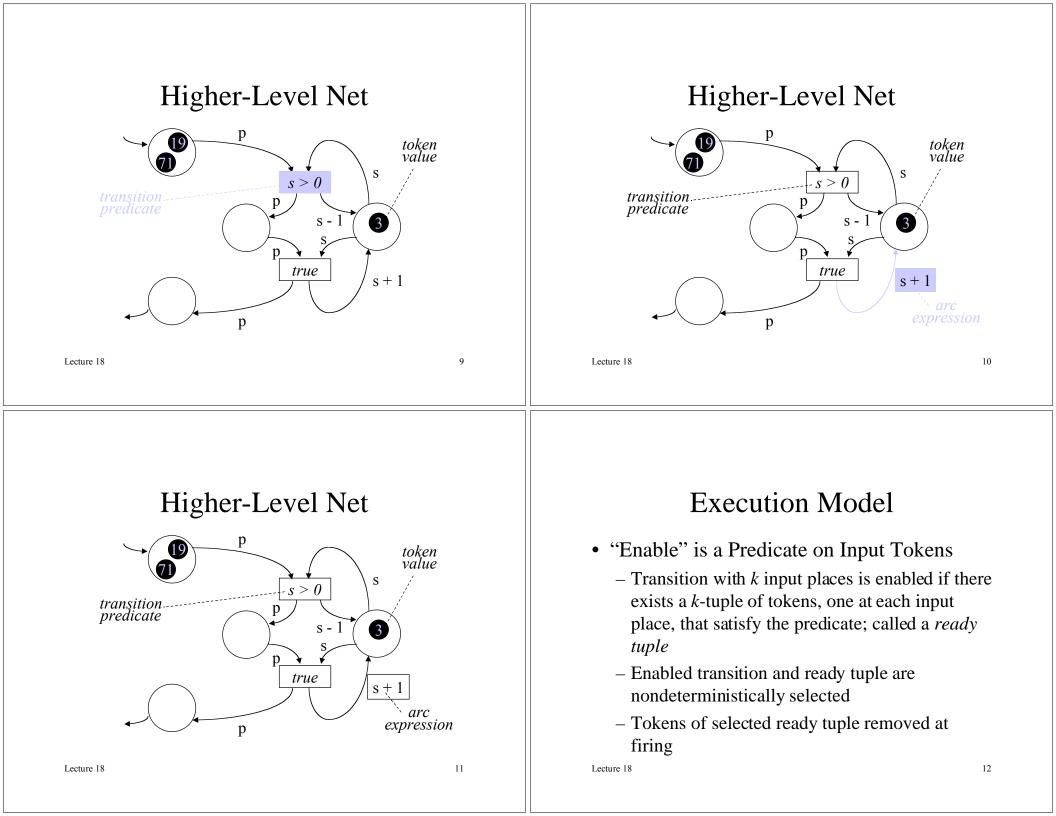
Higher-Level Net



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Execution Model

- Function Computes Output Token Values
 - Transition with *h* output places uses the function to compute *h* values, one for each output token

Higher-Level Net Semaphore

