

Lecture 15

Control Dependence Graphs

Kenneth M. Anderson
CSCI 5828, Spring 2000
This lecture comes from...

A Compositional Approach to Interprocedural Control Dependence Analysis

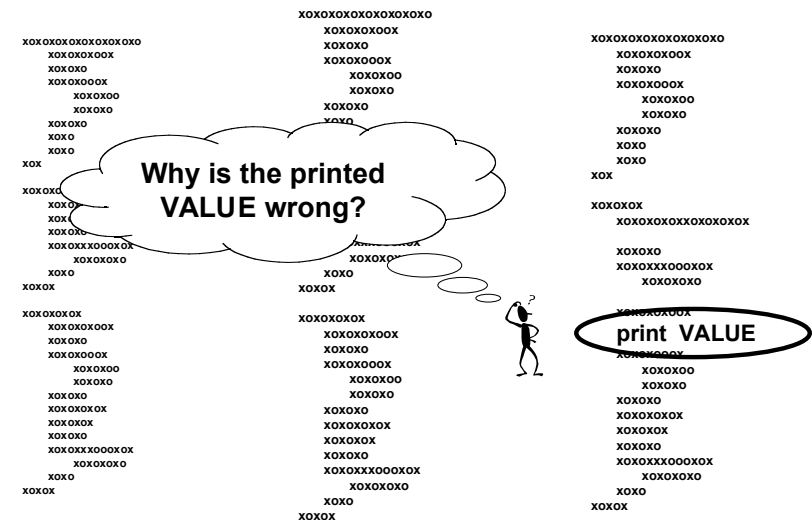
Judith Stafford

Software Engineering Research Laboratory
Department of Computer Science
University of Colorado at Boulder
<http://www.cs.colorado.edu/serl>

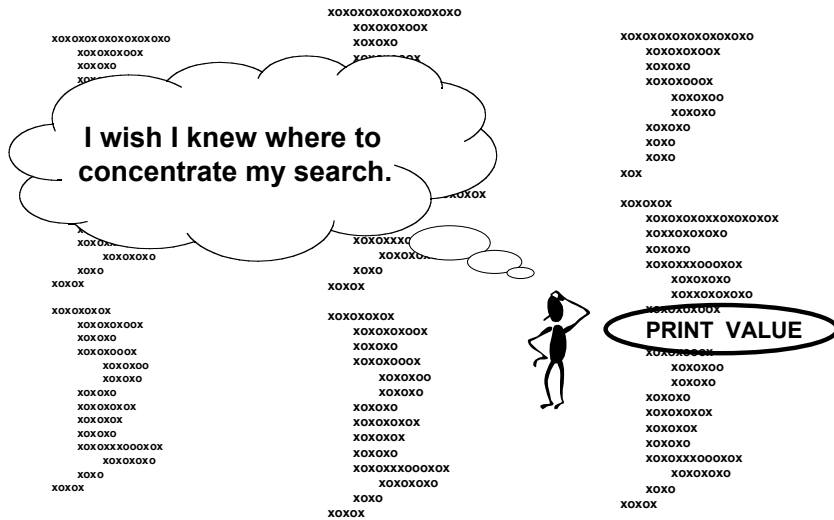
The Roadmap

- Introduction to Dependence Analysis
- ◆ Current State of Affairs and Limitations
- ◆ Judy's Approach -- A Compositional Model
- ◆ Related Work

My Big Program Doesn't Work



But Where To Look?



Simple Bug Tracking Example

- ◆ Question: *What statements of Simple could contain the bug that causes it to always print "1"?*
 - Answer: 1, 2, or 4
 - How do we find the answer?

```
Program Simple
1: read i
2: if ( i == 1)
3:  print "POS:"
   else
4:  i = 1
5:  print i
6: end
```

Getting Started

- ◆ Best to start your search where the bad value is printed -- statement 5
- ◆ It looks like the value used at statement 5 comes from statement 4

```
Program Simple
1: read i
2: if ( i == 1)
3:  print "POS:"
   else
4:  i = 1
5:  print i
6: end
```

Conditional Execution

- ◆ But then you notice that statement 4 might not even be executed because it depends on the decision made at statement 2

```
Program Simple
1: read i
2: if ( i == 1)
3:  print "POS:"
   else
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Variable Assignment

- ◆ The decision made at statement 2 depends on what value is input at statement 1
- ◆ The value printed at 5 may come directly from statement 1

Program Simple

```
1: read i
2: if ( i == 1)
3:  print "POS:"
   else
4:  i = 1
5: print i
6: end
```

The Answer

- ◆ So only statements 1, 2, and 4 could contain the bug...
- * This is helpful

Program Simple

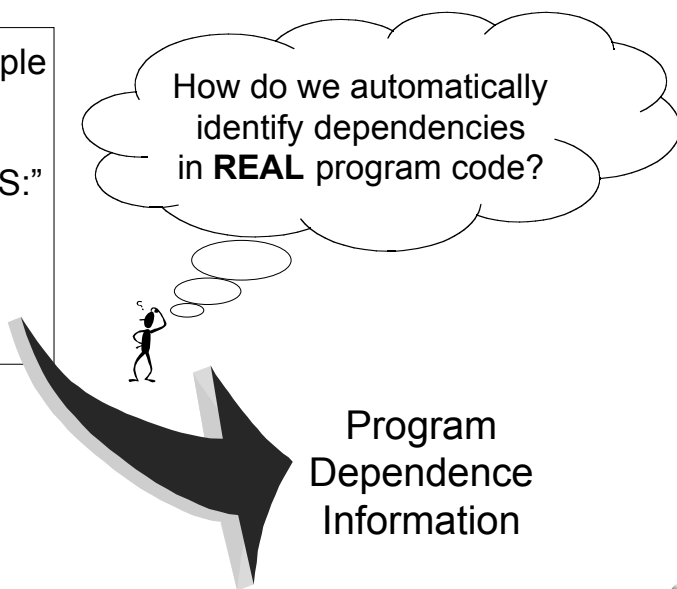
```
1: read i
2: if ( i == 1)
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6: end
```

The Big Question

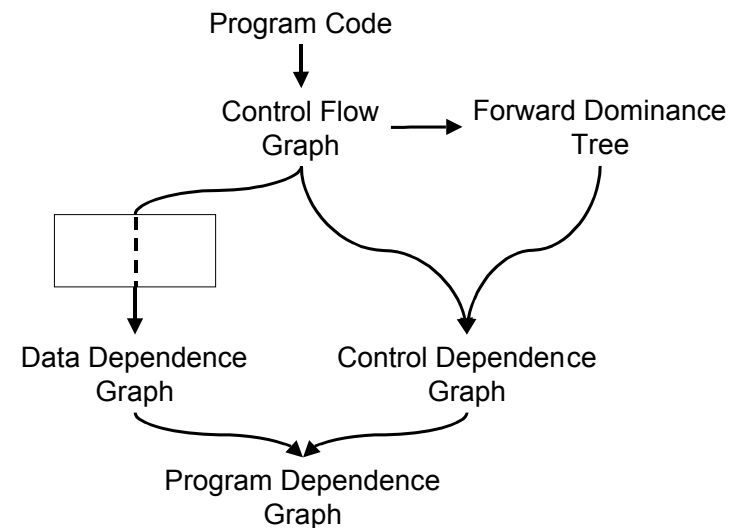
Program Simple

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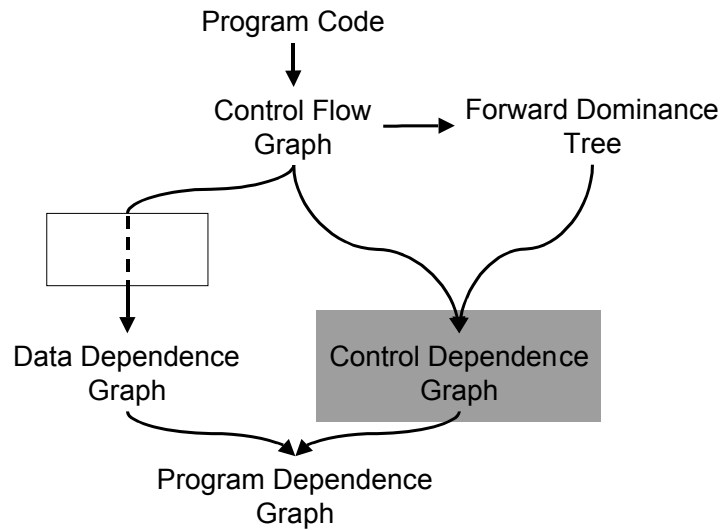
How do we automatically identify dependencies in **REAL** program code?



A Graph-Based Model



A Graph-Based Model



A Control Dependence Representation

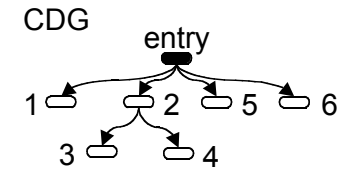
- ◆ Represent control dependencies in a *control dependence graph*, "CDG"

Program Simple

```

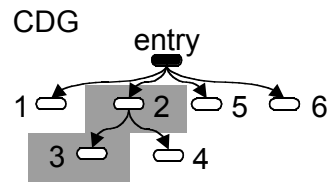
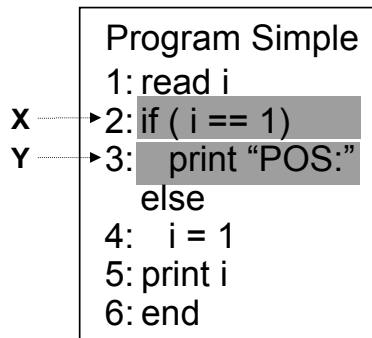
1: read i
2: if ( i == 1 )
3:   print "POS:"
   else
4:   i = 1
5:   print i
6: end
  
```

- Vertices represent executable statements
- Arcs represent direct control dependence
- A distinguished entry vertex representing the start of the program



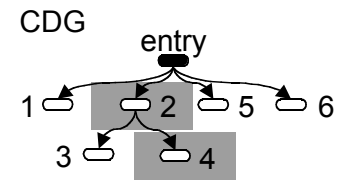
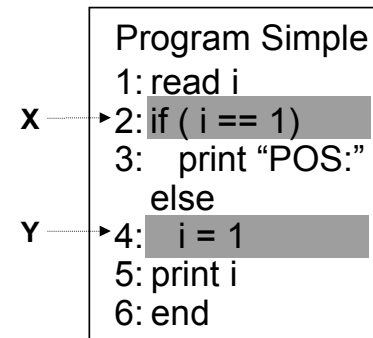
A Control Dependence Representation

- ◆ If statement X determines whether statement Y is executed, statement Y is *control dependent* on statement X



A Control Dependence Representation

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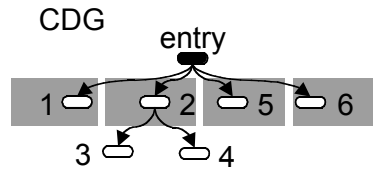


A Control Dependence Representation

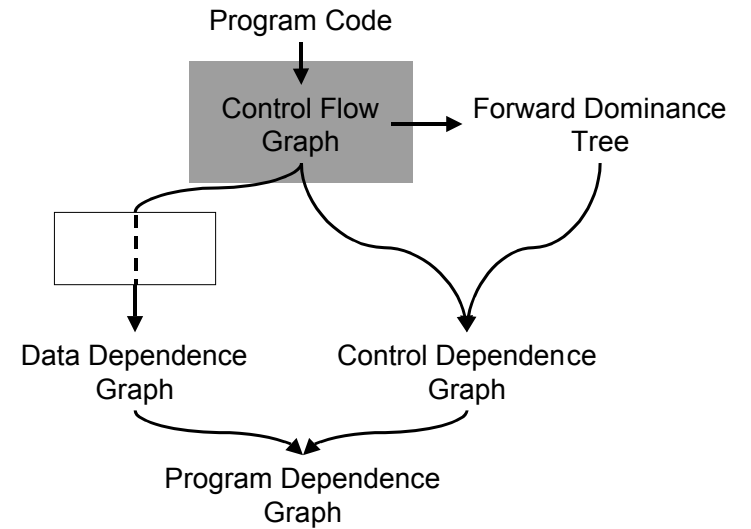
- Statements that are guaranteed to execute are control dependent on entry to the program

```

Program Simple
1: read i
2: if ( i == 1 )
3:  print "POS:"
   else
4:  i = 1
5:  print i
6: end
    
```



A Graph-Based Model



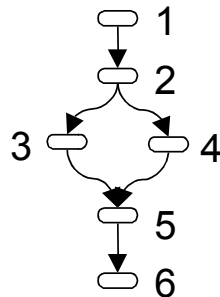
A Graph Representation of Behavior

- The *control flow graph*, "CFG"

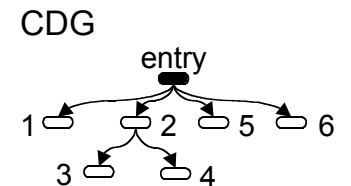
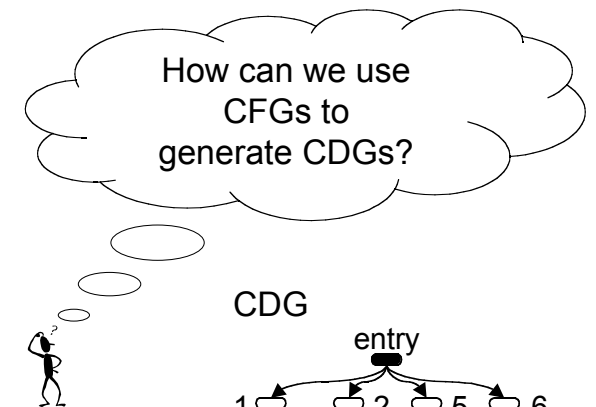
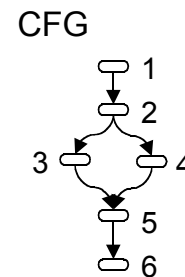
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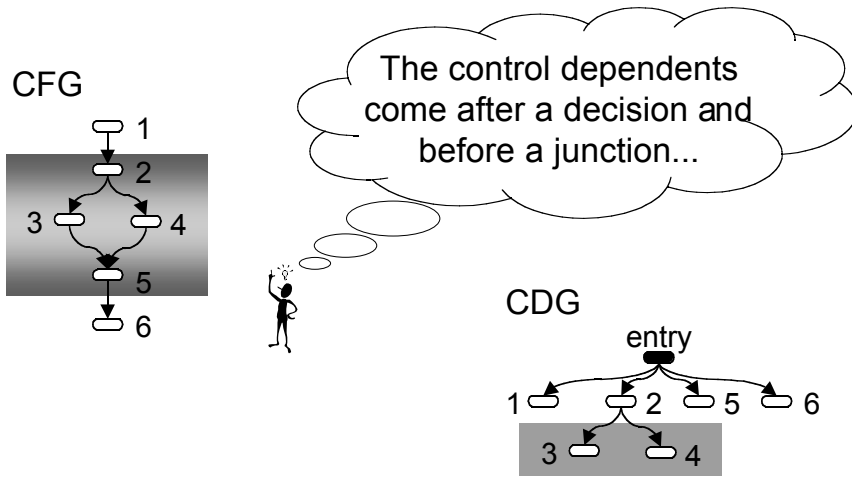
- Vertices represent executable statements
- One entry and one exit
- Arcs represent potential control flow



Calculating Control Dependencies

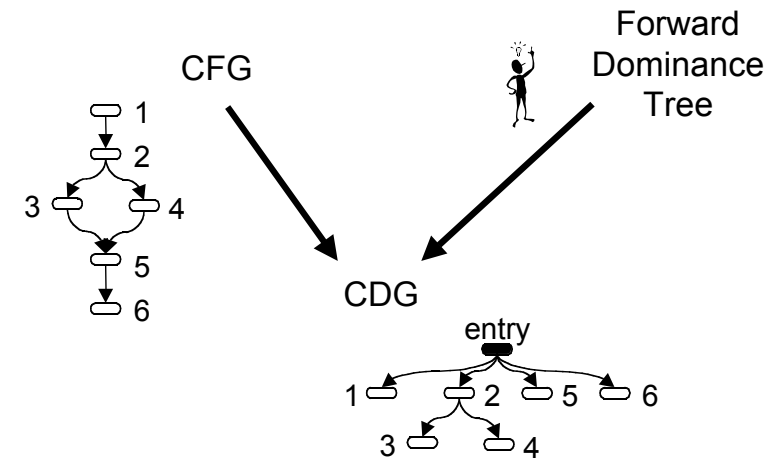


Calculating Control Dependencies

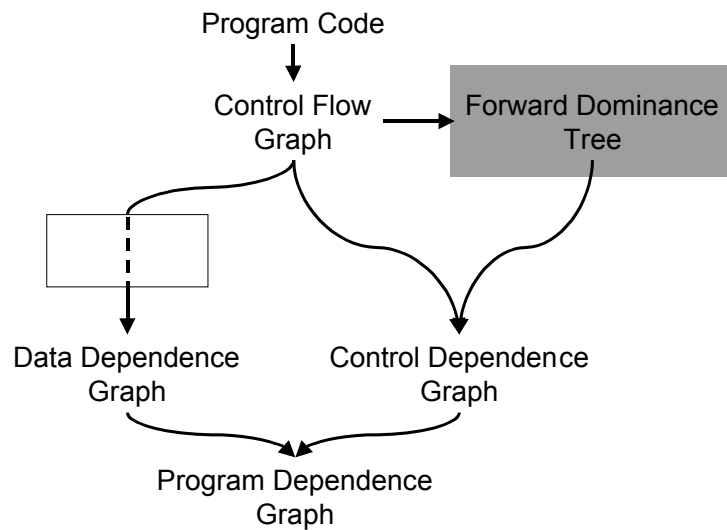


Calculating Control Dependencies

- Use the dominance tree (in reverse)!



A Graph-Based Model



Forward Dominance (a.k.a. post dominance, inverse dominance)

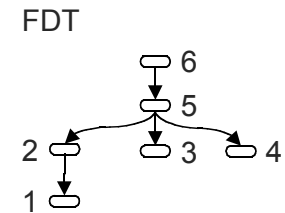
- The *forward dominance tree*, "FDT"

Program Simple

```

1: read i
2: if ( i == 1)
3:   print "POS:"
   else
4:   i = 1
5: print i
6: end
    
```

- Vertices represent executable statements
- Arcs represent immediate forward dominance
- The root of the tree is the exit of the CFG



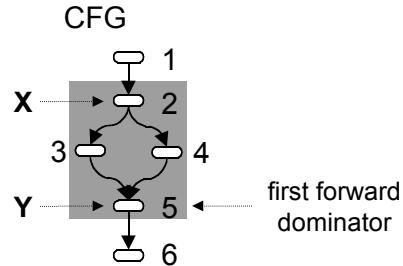
Forward Dominance (a.k.a. post dominance, inverse dominance)

- ◆ Y *forward dominates* X if all paths from X include Y

Program Simple

```

1: read i
2: if ( i == 1)
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   else
4:  i = 1
5:  print i
6: end
    
```



Notice that the control dependents, 3 and 4, don't forward dominate 2

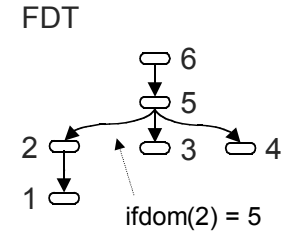
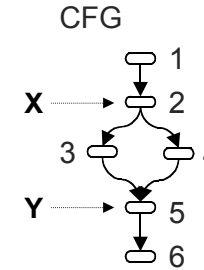
Forward Dominance Tree

- ◆ The first forward dominator of X is called the *immediate forward dominator* of X , " $\text{ifdom}(X)$ "
- ◆ Vertices between X and $\text{ifdom}(X)$ are dependent on X
- ◆ Immediate forward dominators form a tree, "FDT"

Program Simple

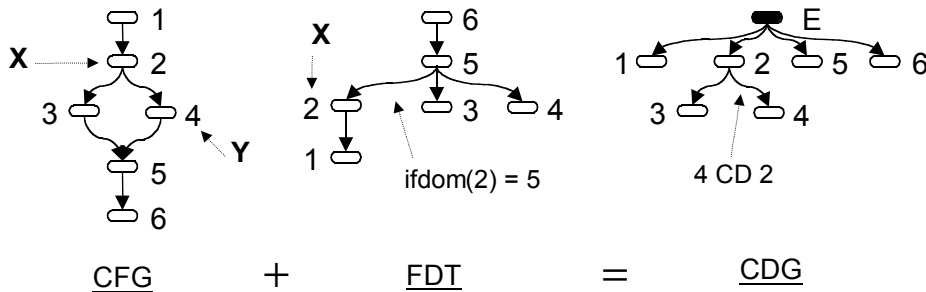
```

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```



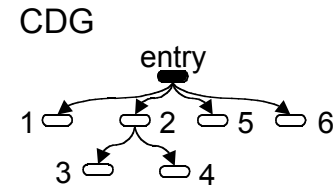
A Graph-Based Definition

- ◆ Y is *control dependent* on $X \Leftrightarrow$ there is a *path in the CFG* from X to Y that doesn't contain the *immediate forward dominator* of X



How Does This Help?

- ◆ Now we have half of the answer



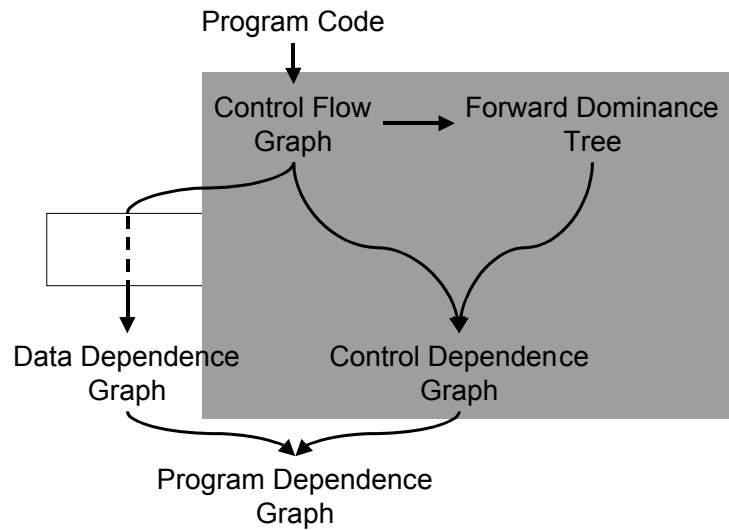
"But then you notice that statement 4 might not even be executed because it depends on the decision made at statement 2"

Program Simple

```

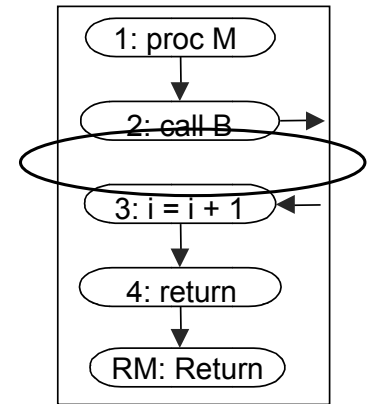
1: read i
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```

A Graph-Based Model



Real Programs are More Complex

- ◆ CFG-based definitions and algorithms expect a connected graph
- ◆ Procedure-level control flow graphs are not connected because there is no direct flow from a call to the next statement



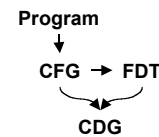
The Roadmap

- ✓ Introduction to Dependence Analysis
- Current State of Affairs and Limitations
- ◆ Judy's Approach -- A Compositional Model
- ◆ Related Work

Other Models

Uni-Procedure Model

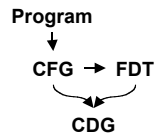
Multi-Procedure Approaches



LIMITED

Other Models

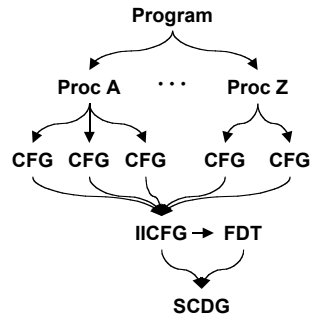
Uni-Procedure Model



LIMITED

Multi-Procedure Approaches

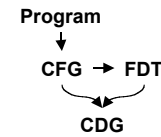
In-lined Approach



IMPRACTICAL

Other Models

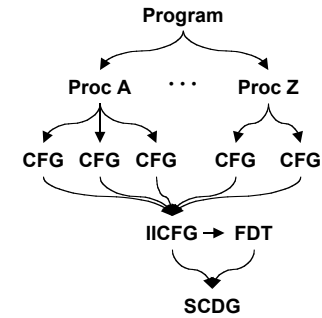
Uni-Procedure Model



LIMITED

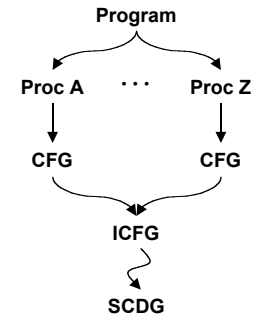
Multi-Procedure Approaches

In-lined Approach



IMPRACTICAL

1-1 Graph Approach



AD HOC

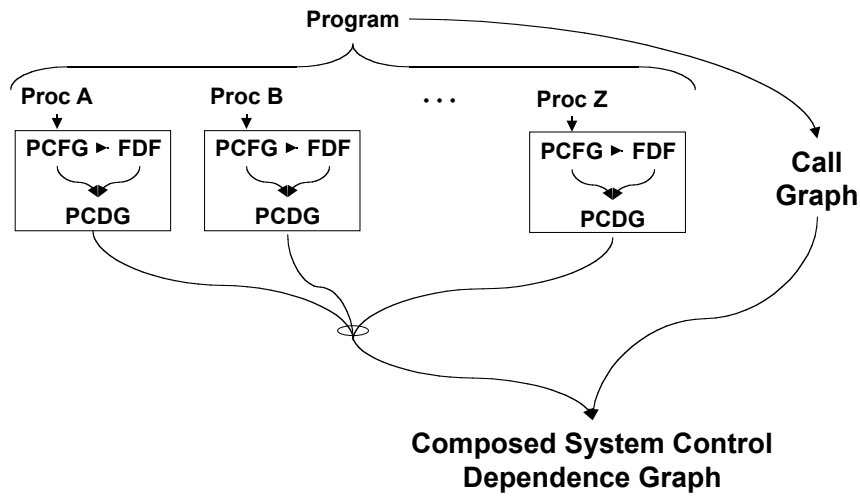
The Roadmap

- ✓ Introduction to Dependence Analysis
- ✓ Current State of Affairs and Limitations
- My Approach -- A Compositional Model
- ◆ Related Work

Guiding Principles

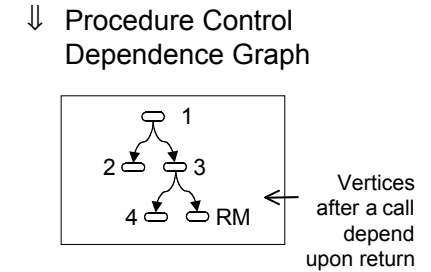
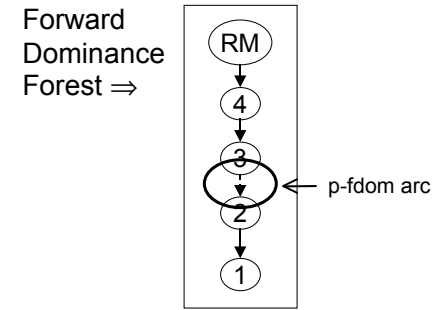
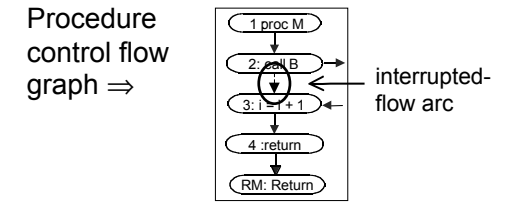
- ◆ Question...
 - Can I extend the forward dominator relation to create a practical and straight-forward model of control dependencies that addresses the pitfalls?
- ◆ Approach
 - Compositional
 - » Reason about properties of procedures independently
 - » Compose procedure-based representations to reflect program-wide properties
 - Language-independent
 - » Modern programs are composed of parts written in different languages
 - Generalizable
 - » Limitations and power are precisely defined

A New View - A Compositional Model



Procedure-level Structures

- Keep track of the potential indirect flows and forward dominances



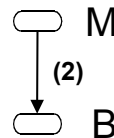
A Graph Representation of Structure

- The *call graph*

Program Multi

- Proc M
- call B
- $i = i + 1$
- return

- Vertices represent procedures in a program
- Arcs represent procedure call
- Arcs are annotated with ID of each call site

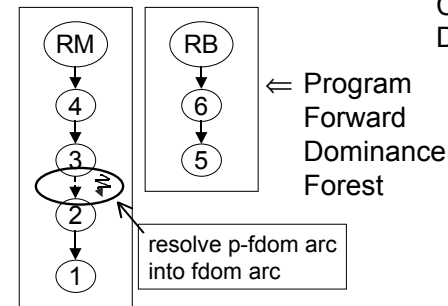
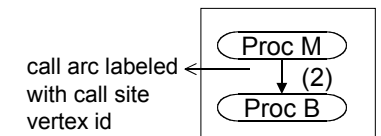


- Proc B
- return

Program-level Structures

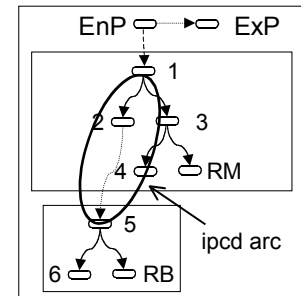
- Apply program call graph to resolve p-fdom arcs and identify interprocedural dependencies

Program Call Graph ↓

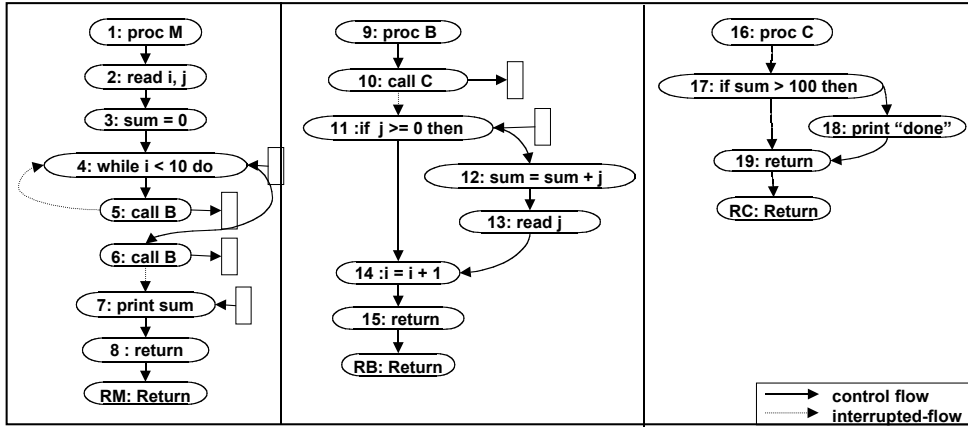


Compound Control Dependence ⇒

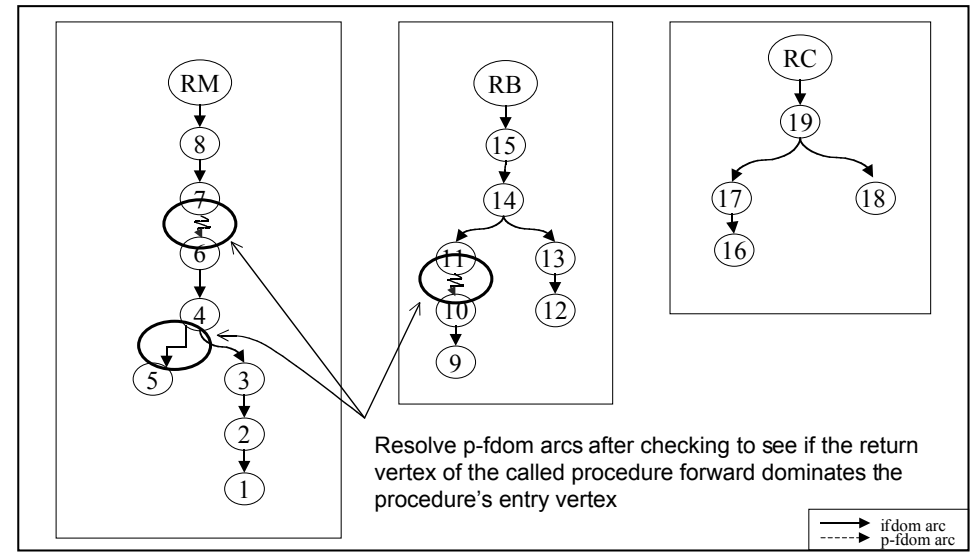
Procedure B and forward dominators of 2 inherit control dependence from 2



Example -- Program Sum



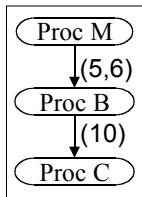
Forward Dominator Forest



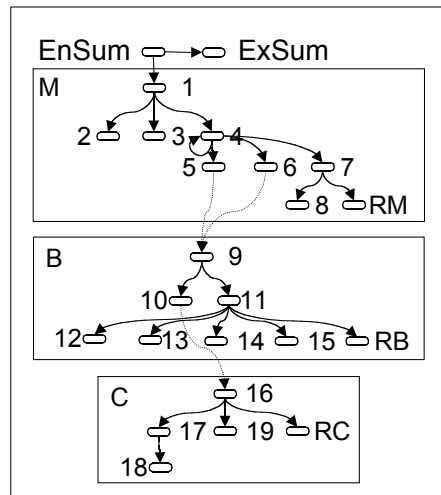
Resolve p-fdom arcs after checking to see if the return vertex of the called procedure forward dominates the procedure's entry vertex

Control Dependencies in Program Sum

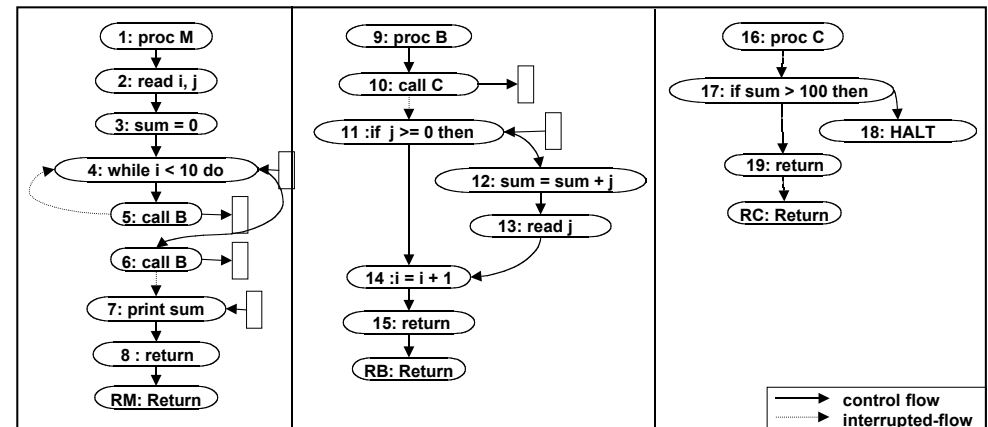
- ◆ Call Graph for Program Sum
- ◆ Sum's Compound Control Dependence Graph



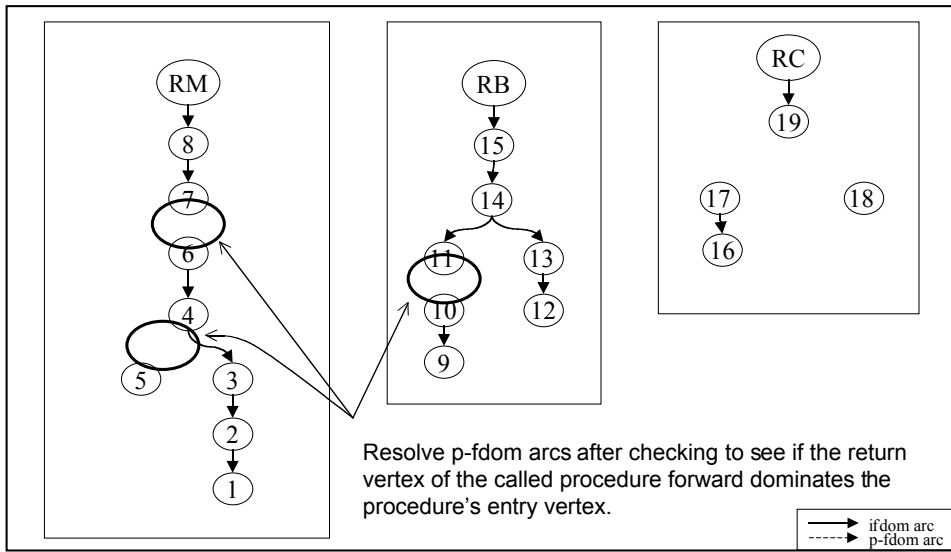
Procedure inherits control dependence of call



Introducing an Embedded Halt

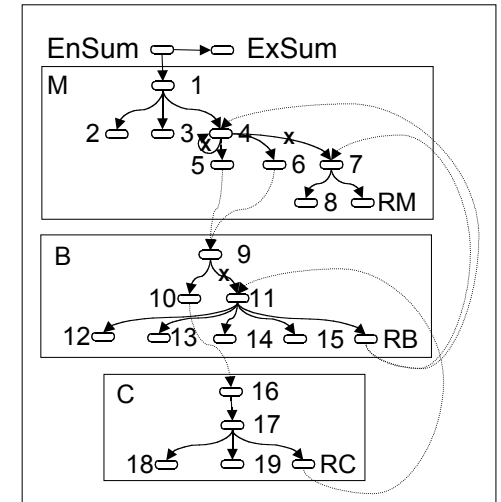
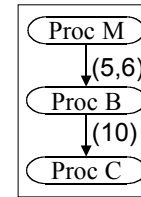


Forward Dominator Forest -- with EHalt



Effect of Halt on Control Dependence

- ◆ Call Graph for Program Sum
- ◆ Sum's Compound Control Dependence Graph



Procedure inherits control dependence of call/return sites

Related Work

- ◆ Researchers have extended the CFG and generate the CDG in *ad hoc* ways to apply to complex programs

CFG + FDT = CDG	Uni-procedure	Podgurski+'90 Ferrante+'87
xFG \rightsquigarrow xCDG	Multi-procedure	Horwitz+'90, Loyall+'93 Harrold+'98,99, Liao+'99
xFG \rightsquigarrow xCDG	Object-oriented	Larsen+'96 Zhao+'96
xFG \rightsquigarrow xCDG	Concurrent	Zhao+'96
xFG \rightsquigarrow xCDG	Concurrent-OO	Hatcliff+'99, Zhao+'99
xFG \rightsquigarrow xCDG	Reactive	Clarke+'99, Stafford+'98