

Lecture 4: Formal SE

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Today's Lecture

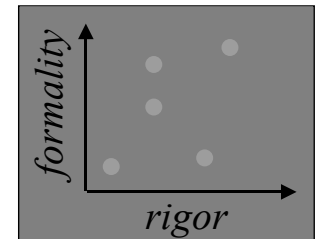
- Introduction to Formal Software Engineering
 - Discuss Models
 - Discuss Formal Notations

Formal Software Engineering

- Software
 - Computer programs and their related artifacts
- Engineering
 - The application of scientific principles in the context of practical constraints
- Formal
 - The use of models, techniques, and tools that are grounded in mathematics

Some Important Points

- *Formal* does not mean *Hard*
- *Formal* does not mean *Good*
- *Informal* does not mean *Bad*
 - unless it means *ad hoc*



What Are “Formal Methods”?

- ❶ *Writing* a formal specification
- ❷ *Proving* properties about the specification
- ❸ *Constructing* a program by mathematically manipulating the specification
- ❹ *Verifying* the program by mathematical argument

Formal SE is Broader

Not just specification and verification of programs...

- Architecture
- Analysis/Testing
- Reliability and Performance Engineering
- Configuration Management
- Process Management, etc.

Model/Specification/Formalism

- Model
 - An abstract representation
- Specification
 - A formal expression of a model or of a property of a model
- Formalism
 - A mathematical notation for writing specifications; a specification language

Specification and the Life Cycle

- Requirements
- Design
 - High level and Low level
- Implementation
- Test

Specification is used in All Activities

Specification/Modeling Styles

- Operational
- Declarative
 - Axiomatic
 - Algebraic
- Structural/Relational

Choice of style dictated by focus of concerns

Specification/Modeling Styles

- Operational (or Imperative)
 - Described according to desired *actions*
 - Usually given in terms of an *execution model*
- Descriptive (or Declarative)
 - Described according to desired *properties*
 - Usually given in terms of *axioms* or *algebras*
- Structural (or Relational)
 - Described according to desired *relationships*
 - Usually given in terms of *multi/hyper graphs*

Logical Foundations

- Predicate/Propositional Logic
- Temporal Logic Systems
- Lambda calculus, etc.

Propositional Logic

- A proposition is a statement that is either true or false, but not both
- Propositional Logic is the language of propositions
 - It consists of well-formed formulas constructed from atomic formulas and logical connectives
 - The meaning of a proposition is determined by the truth values assigned to its assertions

Example

- P = “program does not terminate”
- Q = “alarm rings forever”
- $P \Rightarrow Q$ (If the program does not terminate then alarm rings forever)

• P	Q	$P \Rightarrow Q$
• T	T	T
• T	F	F
• F	T/F	T

Library Example

- S: a book is on the stacks
- R: a book is on reserve
- L: a book is on loan
- Q: a book is requested
- Constraints
 - A book can be in only one of three states S, R, and L
 - If a book is on the stacks or on reserve then it can be requested

Library Example, continued

- Constraints specified as propositions
 - A: $S \Leftrightarrow \neg(R \vee L)$
 - B: $R \Leftrightarrow \neg(S \vee L)$
 - C: $L \Leftrightarrow \neg(S \vee R)$
 - D: $Q \Rightarrow (S \vee R)$
- Prove “if a book is on loan then it is not requested” is a logical consequence

(One Possible) Solution

- Proof by Contradiction
- Steps
 - ① $L \Rightarrow \neg Q$
 - This is our goal
 - ② $\neg(\neg L \vee \neg Q)$
 - We negate it...
 - ③ $L \wedge Q$
 - ...and get this
 - ④ L
 - Conjunction Elim, 3
 - ⑤ $\neg(S \vee R)$
 - Biconditional Elim, 4, C
 - ⑥ Q
 - Conjunction Elim, 3
 - ⑦ $(S \vee R)$
 - Modus Ponens, 6, D
 - ⑧ Contradiction
 - 5 and 7 contradict

Another Solution

- Direct Proof
 - $Q \rightarrow (S \vee R)$
 - $\neg Q \vee (S \vee R)$
 - $L \Leftrightarrow \neg(S \vee R)$
 - $\neg L \Leftrightarrow (S \vee R)$
 - $\neg Q \vee \neg L$
 - $\neg L \vee \neg Q$
 - $L \rightarrow \neg Q$
- Steps
 - D
 - Logical Equivalence of 1
 - C
 - Double Negation of 3
 - Substitution, 2, 4
 - \vee is communicative
 - Logical Equivalence of 6

Predicate Logic

- Propositional Logic cannot specify the relationships between objects
 - It can only assert that particular properties hold or do not hold within a set of propositions
- Predicate Logic has the power to do so
 - consists of
 - constants, predicates, variables, and functions

Example use of predicate logic

- Consider lines and points on a plane
 - (1) two lines meet at a unique point
 - (2) there is a unique line through any two points
 - $\text{line}(x) = x$ is a line
 - $\text{point}(x) = x$ is a point
 - $\text{lies_on}(x, y) = \text{point } x \text{ is contained in line } y$

Example, continued

- domain distinction
 - (a) $\forall x \cdot (\text{point}(x) \vee \text{line}(x));$
 - (b) $\forall x \cdot (\neg(\text{point}(x) \wedge \text{line}(x)));$
- incidence
 - $\forall x, y \cdot (\text{lies_on}(x, y) \Rightarrow (\text{point}(x) \wedge \text{line}(y)));$

Mathematical Foundations

- Set theory
- Graph theory
- Automata theory
- Abstract algebra
- Probability and statistics

Analysis of Specifications

- Static Analysis
Examines specification text to reveal properties
- Dynamic Analysis
Executes specification text to reveal properties

***Choice of analysis dictated by focus of concerns
and choice of specification style***