

Foundations of Network and Computer Security

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Lecture #10
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Announcements

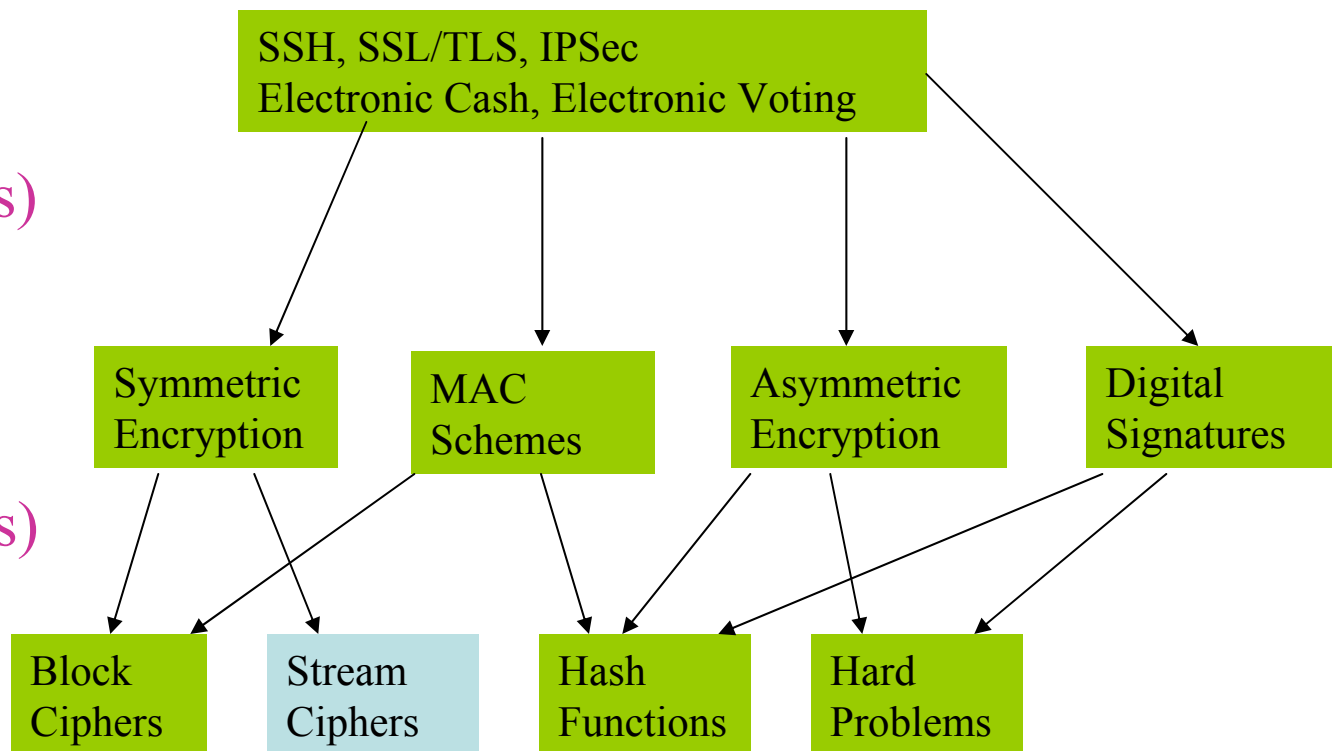
- Quiz #2 next meeting
 - Will cover material up to **last** class (ie, Tuesday's lecture)
 - Up to OpenSSL stuff
 - Nothing from today or next Tuesday's lectures
 - No, you don't have to memorize OpenSSL commands, flags, and oddities
 - Yes, you should know *why* things are done the way they are
 - Review of the RSA stuff on our web site
 - On the course schedule for last lecture
 - There's more than we talked about in class

The Big (Partial) Picture

Second-Level
Protocols
(Can do proofs)

First-Level
Protocols
(Can do proofs)

Primitives
(No one knows how to prove security; make assumptions)



(No one knows how to prove security; make assumptions)

Network Security

- Haven't we already been talking about network security?!
 - Kind of... cryptography is a central part of it
 - Cryptography is nice because it's a neatly packaged science; but we're done for now
 - Network security itself is a vast area with fuzzy borders
 - Research tends to be more ad hoc
 - How do we stop attack A, how do we prevent bug B, how do we detect or tolerate intrusions, etc.

Crypto Good

- The easiest way to break into a computer is usually not by breaking the crypto
 - We've said this a number of times in this class before; there are usually easier ways
- Let's suppose we want to break into a friend's account on CSEL
 - What kind of friend are you??
 - Ok, give me methods... *simple* methods

Breaking into a “Friend’s” Account

- Digression
 - Before we talk about this, let me introduce the “John Disclaimer”
 - I would like each of you to sign a statement “promising not to be evil”
 - I will hand this out at the end of lecture
 - Please remind me
 - It’s also on our web site...
 - Distance students, please print this out and send it in

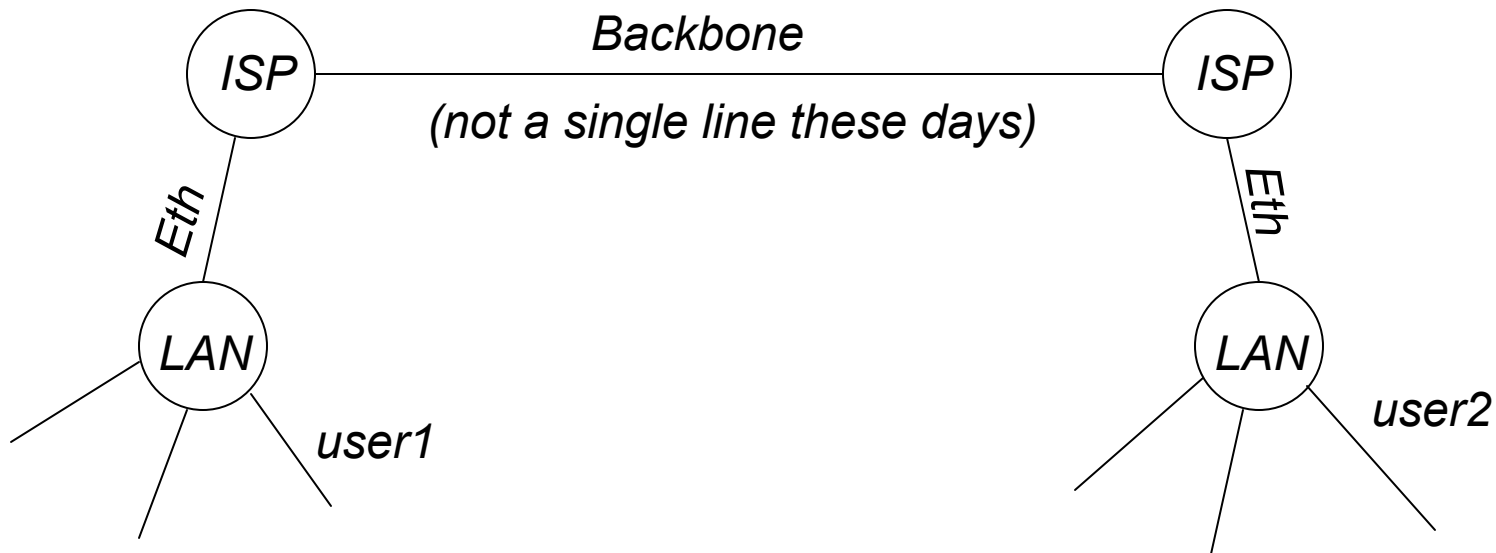
Ok, Breaking into a “Friend’s” Acct

- Fake Login Screen
- Shoulder Surfing
- Password Cracker
 - MD5 hashes publicly available on web
- Social engineering
 - Hard to trick CSOps though
 - Might be easy to impersonate CSOps! 😊
- Key loggers
 - Software and hardware versions
- Keystroke analysis
 - Ok, getting obscure



Networking Refresher

- For some of you this will be boring... sorry
- The basic model:



Basic Networking

- Suppose user1 sends a UDP packet to user2, what happens?
 - What's UDP?
 - User Datagram Protocol
 - Just like IP but with ports
 - Well, first we need an IP address!
 - What's an IP address
 - For IPv4, it's a "dotted quad" of bytes
 - Ex, 128.138.242.21
 - 32 bits
 - For IPv6, it's 128 bits
 - 16 bytes in hex separated by colons

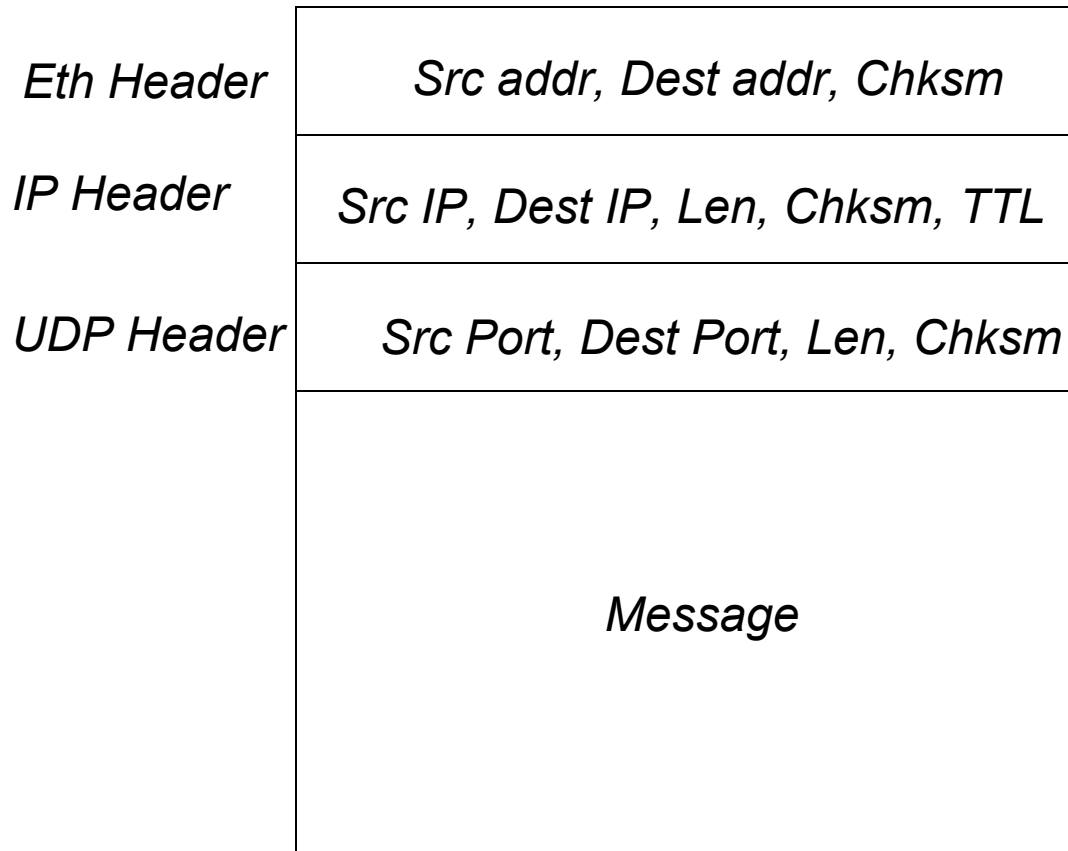
Running out of IP addresses

- 2^{32} is a lot, but we're having problems
 - A lot of hosts out there
 - The class A, B, C scheme is wasteful
 - Though subnetting helps
 - A lot of NAT Boxing “helps”
 - Since we're getting by, it means a slower migration to IPv6

Sending a UDP packet

- Assume IPv4
 - Get IP address via DNS
 - Domain Name Service
 - Distributed database mapping textual names to IP addresses
 - Insecure
 - DNS spoofing
 - More on this later
 - Ok, so we have an IP address
 - And we presumably have a port #

Pack it Up!



Ethernet addresses are called “MAC addresses”

Ethernet checksum is actually appended to end of packet

Ethernet MTU is 1500 bytes

Routing on a Network

- Usually done via OSPF or LSP for LANs
 - Open Shortest Path First, Link-State Protocol
 - These protocols assume “modest sized” networks
 - A routing protocol decides how to forward packets based on routing tables
- BGP is used on backbone
 - Border Gateway Protocol
 - Routes using incomplete information

Local Routing Table

- Our local routing table (on host of user1) is not going to have a route to IP of user2
 - Routing table will therefore send our packet to the gateway
 - Gateway is the machine/router on the “edge” of the network responsible for processing all incoming/outgoing traffic from/to the LAN
 - NAT boxing, firewalling, and other stuff is usually done here as well

Getting to the Gateway

- How to we route to the IP address of the gateway on our local Ethernet?
 - ARP (Address Resolution Protocol)
 - Translates IP addresses into MAC addresses
 - Caches old lookups, so we probably already have the MAC address of the gateway
 - If not, we send an ARP Request to the LAN, including the IP address whose MAC we seek
 - Owner (ie, the gateway) sends ARP Reply with his MAC address and we cache it
 - Usually, all other machines who hear the ARP Reply cache it as well
 - Leads to attacks... more later

Sending to the Gateway

- Now we have the MAC address of the gateway
 - Send our packet to the gateway via the Ethernet protocol
 - This is usually done with a hardware device (network card) which often puts the Eth header on your packet for you, computes checksums, etc.
 - Broadcasts packet, detects collisions
 - Exponential backoff
 - Promiscuous mode – Sniffers use this
 - Works through hubs, but doesn't work through switches on a switched Ethernet
 - You can often fool switches

Gateway Receives Eth Packet

- Strips Eth header and again tries to route the resulting IP packet
 - Looks in routing table, sends to ISP
 - ISP probably routes using BGP
 - Reaches other ISP
 - Note that we're using other Ethernets and similar physical-layer protocols for each hop!
 - Other ISP routes to other LAN's gateway
 - Gateway sees IP is in its range and does ARP to route to user2

User2 Receives Packet

- User2 receives the IP packet
 - Removes IP header
 - No one else (is supposed to) look inside packet until user2 receives it
 - NAT boxes break this rule
 - Firewalls break this rule
 - See it's a UDP packet and “sends” to proper port
 - Ports are mapped to applications via `listen()`
 - Application receives message and processes it

Other Protocols

- We didn't even talk about SLIP or PPP
- ATM, FDDI, Wireless
- What about DHCP?
 - Dynamic IP addresses
- There is also ICMP
 - Internet Control Message Protocol
 - Echo (ping), traceroute
- Application Layer Protocols
 - SNMP – Network Management
 - SMTP – Sendmail
 - POP/IMAP – Mail protocols

MTU – Maximum Transmission Unit

- MTU for Ethernet is 1500 bytes
 - If MTU is exceeded, packet is “fragmented”
 - IP has support for packet fragmentation and reassembly
 - A packet is broken into as many pieces as necessary to comply with MTU
 - Fragments routed as regular IP datagrams, independent of each other
 - Reassembly done at host only

IP – Best Effort Datagrams

- IP is “best effort”
 - There is no tracking of packets
 - If something is dropped... oh well
 - If one fragment is dropped, many transport layer protocols (like TCP) will consider the whole thing lost and not ACK
 - This seems bad, but it’s one of the biggest successes of IP
 - UDP is IP with ports, so it too is “best effort”

TCP – Transmission Control Protocol

- Stateful connections
 - Runs over IP just like UDP, but adds more than just ports
 - Establish a connection with `listen()` and `connect()`
 - IP and UDP were “stateless” protocols
 - Reliable delivery
 - Unlike best-effort, this protocol guarantees delivery of packets, in proper order
 - Uses sequence numbers, sliding windows, ACKs every transmission

Crypto on a Network

- How do we do crypto on a network?
 - We've seen application-layer examples
 - SSL/TLS, SSH
 - This is called “end-to-end” cryptography, meaning between hosts
 - The routers don't care if the innermost part of each packet (the “payload”) is ciphertext or plaintext
 - IPSec
 - IPSec does crypto at the network layer (the IP layer)
 - Extremely well-engineered; hardly used
 - We won't study IPSec in this course

Network Security: The Biggest Challenges

- What are the biggest problems now, today, on the Internet
 - What are the most common types of attacks?
 - Viruses, worms
 - Break-ins via software vulnerabilities
 - Denial of Service attacks (DoS)
 - And Distributed Denial of Service (DDoS)
 - What about keyloggers, spyware, rootkits?
 - Not as relevant to network security
 - More likely to be end-results of other break-ins
 - A recent virus was found to install a keylogger

Viruses (Worms)

- Today, most everyone just calls them viruses
 - Technically most are “worms”
 - Worm is a self-contained propagating program
 - Viruses embed in other programs and self-replicate
 - Kind of like viruses in biology

Viruses: History

- Morris Worm, Nov 2nd, 1988
 - The first worm (I know of) was the Morris worm
 - Robert T. Morris, Jr.
 - 23 years old
 - Cornell grad student
 - Father worked at the NSA (whoops!)
 - Wrote a self-propagating program as a “test concept”
 - Exploited Unix vulnerabilities in sendmail and fingerd
 - Released at MIT
 - Bug in the worm caused it to go wild
 - Probably wouldn’t have caused much damage otherwise!

Morris Worm (cont)

- Shut down thousands of Unix hosts
 - But this was 1988...
- Reactions
 - People didn't know what to do, so they panicked
 - Disconnected from net
 - Unable to receive patches!
 - Morris fined \$10k, 3 yrs probation, 400 hrs community service
 - CERT was created

CERT -- They were first

- Carnegie Mellon Emergency Response Team
 - But don't expand it into an acronym
- Provide technical advice and coordinate responses to security compromises
- Identify trends in intruder activity
- Work with other security experts to identify solutions to security problems
- Disseminate information to the broad community
- Analyze product vulnerabilities
- Publishes technical documents
- Presents training courses

Modern Viruses

- Almost all look for Windows hosts
 - Windows runs on more than 90% of desktops these days
 - A lot of hosts on cable modems
 - Fast, always on
 - Destructive payloads
 - Wipe hard disk, eg
 - Some install backdoors for later use
 - All kinds of weird behaviors though
 - Some innocuous

Viruses: Why?

- Who writes these things?
 - Typical profile: male, teenager, geeky, smart
 - Script Kiddies
 - Don't really write them, but launch them
 - Sometimes make small mods and call them their own
 - Scariest hackers: beyond the reach of the law
- Why?
 - Intellectual challenge (sigh...)
 - Peer recognition
 - Bot building (Zombie armies)
 - Because it's there?