

ML Type Inference and Unification

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Research Goals

- Easy to use, high performance parallel programming
 - Primary contributions in backend and runtime
 - Need a front end to target backend
 - ML offers ease of use and safety
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ML Type Inference

- Hindley/Milner Type Inference
- Statically typed language with no mandatory annotations
- Three phases to determining types
 - Constraint generation
 - Unification
 - Annotation



An Example

```
let rec apply = fun f v t ->  
  if t = 0 then  
    v  
  else  
    apply f (f v) (t - 1)  
fi
```



An Example

```
let rec apply = fun f v t ->  
  if t = 0 then  
    v  
  else  
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fi
```

```
val apply: ('a->'a)->'a->int->'a
```

Constraint Generation

```
let rec apply = fun f v t ->  
  if t = 0 then  
    v  
  else  
    apply f (f v) (t - 1)  
fi
```

Variables

Constraints



Constraint Generation

```
let rec apply = fun f v t ->  
  if t = 0 then  
    v  
  else  
    apply f (f v) (t - 1)  
fi
```

Variables

apply: 'a

Constraints



Constraint Generation

```
let rec apply = fun f v t ->  
  if t = 0 then  
    v  
  else  
    apply f (f v) (t - 1)  
fi
```

Variables

apply: 'a
f: 'b
v: 'c
t: 'd

Constraints



Constraint Generation

```
let rec apply = fun f v t ->  
  if t = 0 then  
    v  
  else  
    apply f (f v) (t - 1)  
fi
```

Variables

apply: 'a
f: 'b
v: 'c
t: 'd

Constraints

'a = 'b → 'e

Constraint Generation

```
let rec apply = fun f v t ->  
  if t = 0 then  
    v  
  else  
    apply f (f v) (t - 1)  
fi
```

Variables

apply: 'a
f: 'b
v: 'c
t: 'd

Constraints

'a = 'b → 'e
'b = 'c → 'f

Constraint Generation

```
let rec apply = fun f v t ->  
  if t = 0 then  
    v  
  else  
    apply f (f v) (t - 1)  
fi
```

Variables

apply: 'a
f: 'b
v: 'c
t: 'd

Constraints

'a = 'b → 'e
'b = 'c → 'f
'e = 'f → 'g

Constraint Generation

```
let rec apply = fun f v t ->  
  if t = 0 then  
    v  
  else  
    apply f (f v) (t - 1)  
fi
```

Variables

apply: 'a
f: 'b
v: 'c
t: 'd

Constraints

'a = 'b → 'e
'b = 'c → 'f
'e = 'f → 'g
'd = int

Constraint Generation

```
let rec apply = fun f v t ->  
  if t = 0 then  
    v  
  else  
    apply f (f v) (t - 1)  
fi
```

Variables

apply: 'a
f: 'b
v: 'c
t: 'd

Constraints

'a = 'b → 'e
'b = 'c → 'f
'e = 'f → 'g
'd = int
'g = int → 'h

Constraint Generation

```
let rec apply = fun f v t ->  
  if t = 0 then  
    v  
  else  
    apply f (f v) (t - 1)  
fi
```

Variables

```
apply: 'a  
f: 'b  
v: 'c  
t: 'd
```

Constraints

```
'a = 'b → 'e  
'b = 'c → 'f  
'e = 'f → 'g  
'd = int  
'g = int → 'h  
'd = int  
bool = bool
```

Constraint Generation

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let rec apply = fun f v t ->  
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Variables

```
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f: 'b  
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t: 'd
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Constraints

```
'a = 'b → 'e  
'b = 'c → 'f  
'e = 'f → 'g  
'd = int  
'g = int → 'h  
'd = int  
bool = bool  
'c = 'h
```

Constraint Generation

```
let rec apply = fun f v t ->  
  if t = 0 then  
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fi
```

Variables

apply: 'a
f: 'b
v: 'c
t: 'd

Constraints

'a = 'b → 'e
'b = 'c → 'f
'e = 'f → 'g
'd = int
'g = int → 'h
'd = int
bool = bool
'c = 'h
'a = 'b → 'c → 'd → 'c

Constraint Solving - Unification

Constraints

'a = 'b → 'e

'b = 'c → 'f

'e = 'f → 'g

'd = int

'g = int → 'h

'd = int

bool = bool

'c = 'h

'a = 'b → 'c → 'd → 'c

Mapping



Constraint Solving - Unification

Constraints

'b = 'c → 'f

'e = 'f → 'g

'd = int

'g = int → 'h

'd = int

bool = bool

'c = 'h

'b → 'e = 'b → 'c → 'd → 'c

Mapping

'a = 'b → 'e

Constraint Solving - Unification

Constraints

'e = 'f → 'g

'd = int

'g = int → 'h

'd = int

bool = bool

'c = 'h

('c → 'f) → 'e = ('c → 'f) → 'c → 'd → 'c

Mapping

'a = ('c → 'f) → 'e

'b = 'c → 'f

Constraint Solving - Unification

Constraints

'd = int

'g = int → 'h

'd = int

bool = bool

'c = 'h

('c → 'f) → 'f → 'g = ('c → 'f) → 'c → 'd → 'c

Mapping

'a = ('c → 'f) → 'f → 'g

'b = 'c → 'f

'e = 'f → 'g

Constraint Solving - Unification

Constraints

'g = int \rightarrow 'h

int = int

bool = bool

'c = 'h

('c \rightarrow 'f) \rightarrow 'f \rightarrow 'g = 'c \rightarrow ('f \rightarrow 'c) \rightarrow int \rightarrow 'c

Mapping

'a = ('c \rightarrow 'f) \rightarrow 'f \rightarrow 'g

'b = 'c \rightarrow 'f

'e = 'f \rightarrow 'g

'd = int

Constraint Solving - Unification

Constraints

$\text{int} = \text{int}$

$\text{bool} = \text{bool}$

$'c = 'h$

$('c \rightarrow 'f) \rightarrow 'f \rightarrow \text{int} \rightarrow 'h = ('c \rightarrow 'f) \rightarrow 'c \rightarrow \text{int} \rightarrow 'c$

Mapping

$'a = ('c \rightarrow 'f) \rightarrow 'f \rightarrow \text{int} \rightarrow 'h$

$'b = 'c \rightarrow 'f$

$'e = 'f \rightarrow \text{int} \rightarrow 'h$

$'d = \text{int}$

$'g = \text{int} \rightarrow 'h$

Constraint Solving - Unification

Constraints

'c = 'h

('c → 'f) → 'f → int → 'h = ('c → 'f) → 'c → int → 'c

Mapping

'a = ('c → 'f) → 'f → int → 'h

'b = 'c → 'f

'e = 'f → int → 'h

'd = int

'g = int → 'h

Constraint Solving - Unification

Constraints

$('c \rightarrow 'f) \rightarrow 'f \rightarrow \text{int} \rightarrow 'c = ('c \rightarrow 'f) \rightarrow 'c \rightarrow \text{int} \rightarrow 'c$

Mapping

$'a = ('c \rightarrow 'f) \rightarrow 'f \rightarrow \text{int} \rightarrow 'c$

$'b = 'c \rightarrow 'f$

$'e = 'f \rightarrow \text{int} \rightarrow 'c$

$'d = \text{int}$

$'g = \text{int} \rightarrow 'c$

$'h = 'c$

Constraint Solving - Unification

Constraints

'c → 'f = 'c → 'f

'f = 'c

int = int

'c = 'c

Mapping

'a = ('c → 'f) → 'f → int → 'c

'b = 'c → 'f

'e = 'f → int → 'c

'd = int

'g = int → 'c

'h = 'c

Constraint Solving - Unification

Constraints

'c = 'c

'f = 'f

'f = 'c

int = int

'c = 'c

Mapping

'a = ('c → 'f) → 'f → int → 'c

'b = 'c → 'f

'e = 'f → int → 'c

'd = int

'g = int → 'c

'h = 'c

Constraint Solving - Unification

Constraints

'f = 'c
int = int
'c = 'c

Mapping

'a = ('c → 'f) → 'f → int → 'c
'b = 'c → 'f
'e = 'f → int → 'c
'd = int
'g = int → 'c
'h = 'c

Constraint Solving - Unification

Constraints

int = int

'c = 'c

Mapping

'a = ('c → 'c) → 'c → int → 'c

'b = 'c → 'c

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'd = int

'g = int → 'c

'h = 'c

'f = 'c

Constraint Solving - Unification

Constraints

Mapping

'a = ('c → 'c) → 'c → int → 'c

'b = 'c → 'c

'e = 'c → int → 'c

'd = int

'g = int → 'c

'h = 'c

'f = 'c

Type Annotation

```
let rec apply = fun f v t ->  
  if t = 0 then  
    v  
  else  
    apply f (f v) (t - 1)  
fi
```

Variables

```
apply: 'a  
f: 'b  
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```

Mapping

```
'a = ('c → 'c) → 'c → int → 'c  
'b = 'c → 'c  
'e = 'c → int → 'c  
'd = int  
'g = int → 'c  
'h = 'c  
'f = 'c
```

Type Annotation

```
let rec apply = fun f v t ->  
  if t = 0 then  
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fi
```

Variables

```
apply: ('c → 'c) → 'c → int → 'c  
f: 'c → 'c  
v: 'c  
t: int
```

Mapping

```
'a = ('c → 'c) → 'c → int → 'c  
'b = 'c → 'c  
'e = 'c → int → 'c  
'd = int  
'g = int → 'c  
'h = 'c  
'f = 'c
```

Type Annotation

```
let rec apply : ('c -> 'c) -> 'c -> int -> 'c
  = fun (f:'c -> 'c) (v:'c) (t:int) ->
    if t = 0 then
      v
    else
      apply f (f v) (t - 1)
  fi
```

Variables

apply: ('c → 'c) → 'c → int → 'c

f: 'c → 'c

v: 'c

t: int

Difficulties

- Polymorphic function application
- Matching
- Reference Types



Polymorphic Function Application

```
let f : 'a->'a = fun (x: 'a) -> x
```

```
let t1 = f true
```

```
let t2 = f 3
```

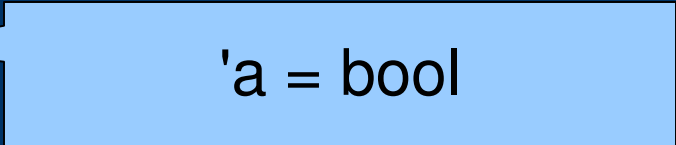


Polymorphic Function Application

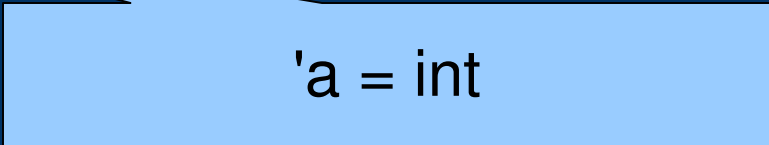
```
let f : 'a->'a = fun (x: 'a) -> x
```

```
let t1 = f true
```

```
let t2 = f 3
```



'a = bool



'a = int

Solution

- Copy the type of f every time f is used

```
let f : 'a->'a = fun (x: 'a) -> x
```

```
let t1 = f true
```

'b = bool

```
let t2 = f 3
```

'c = int



Matching

- Different types for expression being matched and that used with unions:


```
type 'a list =  
  | Nil  
  | Cons of 'a * 'a list
```

```
let map = fun f l ->  
  case l  
  | Nil -> Nil  
  | Cons(h,t) -> Cons(f h, map f t)  
  esac
```

Matching

- Different types for expression being matched and that used with unions:

```
type 'a list =  
  | Nil  
  | Cons of 'a * 'a list  
  
let map = fun f l ->  
  case l  
  | Nil -> Nil  
  | Cons(h,t) -> Cons(f h, map f t)  
  esac
```



l : 'a list

Matching

- Different types for expression being matched and that used with unions:

```
type 'a list =  
  | Nil  
  | Cons of 'a * 'a list
```

```
let map = fun f l ->  
  case l  
  | Nil -> Nil  
  | Cons(h,t) -> Cons(f h, map f t)  
  esac
```

Cons(h,t) :
'a * 'a list

Solution

- Folding and Unfolding
 - λ is folded
 - **Cons**(**h**, **t**) is unfolded
 - Implicit in ML
-
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Reference Types

- Classical ML Bug:

```
let r = ref (fun x -> x)
r := (fun x -> x + 1)
!r true
```



Solution

- Value Restriction
 - SML
 - Only allow values
 - Modified Value Restriction
 - OCaml
 - Value assigned at first use
 - Monomorphic in use, polymorphic at initial definition
-
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Conclusion

- In restricted type systems, full inference can be performed through unification
 - Allows code compactness and static type safety
- Type rules contain constraint generation
- Unification uses constraints to reduce potential solutions to the one correct one



References

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