## Quantum

## Communication and Cryptography

## PHYS/CSCl 3090

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## Come see us!

- Alexandra Kolla/ Graeme Smith: Friday 3:00-4:00 pm, JILA X3I7.
- Ariel Shlosberg:Tu/Th 2:00-4:00pm, DUANG2B90 (physics help room)
- Steven Kordonowy:Th I Iam-I2pm, ECAE 124.
- Matteo Wilczak:Wednesday, I-2pm, DUANG2B90 (physics help room)

New topic: few qubit protocols
(Chapter 6)

- Looking at some simple protocols on just a few qubits.
- Not exactly computing, but more about communication (and also cryptography)
- Last class: started BB84 quantum key distribution protocol.


## Alice, Bob, and Eve

My account number is 330495 -19A.


- Alice wants to send a secret message / binary string to Bob
- She cannot risk Eve learning it


## One Time Pad



- Alice wants to send a secret message / binary string
- She cannot risk Eve learning it
- If Alice XORS her message $x$ with a random $n$-bit string, then she sends a random message and eve cannot learn anything!
- If Bob knew $r$, he could decode $x$.
- How can they share a random string $r$, that Eve cannot find?


## One Time Pad: how to get one

- Make two copies of a bunch of random strings.
- Use a trusted courier to send the pad to the person you want to talk to
- Put it in a locked briefcase on the way.
- Use a trusted courier to send the key.
- But. What if the courier is a bad guy.


## One Time Pads



Actual Pad


Leo Marks, Special Operations Executive

One approach: everyone carries around a bunch of random bits. One "key" for each person you need to talk to.

Not great:
I) you need lots of different pads.
2) someone can peek at your pad, then learn your messages. Have to guard the pad constantly!

## One Time Pad: how to get one

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- Use a trusted courier to send the pad to the person you want to talk to
- Put it in a locked briefcase on the way.
- Use a trusted courier to send the key.
- But. What if the courier is a bad guy.
- With BB84, you don't have to trust anybody. Trust quantum mechanics instead!


## BB84 (Bennett-Brassard-I984)

Alice sends Bob a long sequence of qubits randomly chosen to be in one of the four states:
$|0\rangle,|1\rangle$ (type 1)
$H|0\rangle=1 / \sqrt{2}(|0\rangle+|1\rangle), H|1\rangle=1 / \sqrt{2}(|0\rangle-|1\rangle)$
(type H)
When Alice wants to send the bit 0 , she randomly sends either
$|0\rangle$ or $\mathrm{H}|0\rangle$.
-When she wants to send the bit I, she randomly sends either $|1\rangle$ or $\mathrm{H}|1\rangle$.

- Bob, once he receives each qubit, he randomly decides to apply either I (type I measurement) or H (type H measurement) to the qubit and then measure in the standard basis.


## The protocol

$$
x=0011
$$

Allice


Bob


- If both Alice and Bob chose type I or they both chose type H, then their respective bits agree after Bob measures.
- How do they know which bits agree?


## The protocol



- If both Alice and Bob chose type I or they both chose type H, then their respective bits agree after Bob measures.
- How do they know which bits agree?
- Alice sends Bob over an insecure channel, which qubits were type I and type H .
- She does not reveal if the bit was 0 or I to begin with.
- For those qubits that Alice's choice agrees with Bob's measurement, Bob learns the actual value of the original bit Alice wanted to send.
- They throw out the rest.


## A possible attack by Eve

- During transmission, Alice sends one of four states: $|0\rangle,|1\rangle,|+\rangle,|-\rangle$
- Eve wants to know which one is sent
- A natural attack is for Eve to pick a random basis and measure in that.
- Let's say Alice transmits in $\{|0\rangle,|1\rangle\}$, Eve Measures $\{|0\rangle,|1\rangle\}$, and Bob measures $\{|0\rangle,|1\rangle\}$.
- Eve and Bob both learn the bit. Not good!


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-What then?


## Concept Test

Alice transmits $|1\rangle$,
Eve Measures $\{|+\rangle,|-\rangle\}$
Bob measures $\{|0\rangle,|1\rangle\}$.
What is the probability Bob measures |1〉?
A) 0
B) 1
C) 0.5
D) 0.75

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## A possible attack by Eve

- During transmission, Alice sends one of four states: $|0\rangle,|1\rangle,|+\rangle,|-\rangle$
- Eve wants to know which one is sent
- A natural attack is for Eve to pick a random basis and measure in that.
- If she does this, it will introduce some disagreements in Alice's and Bob's bit strings.
- By sacrificing a small fraction of their key (publicly announcing the values and comparing) they can detect such meddling.
- If their keys are always in agreement, then they can have high confidence no eavesdropping has occurred.


## A perfect attack!

- What eve really wants is an operation that does this:
- $|0\rangle \rightarrow|0\rangle|0\rangle$
- $|1\rangle \rightarrow|1\rangle|1\rangle$
- $|+\rangle \rightarrow|+\rangle|+\rangle$
- $|-\rangle \rightarrow|-\rangle|-\rangle$
- If she had such a machine, she could learn everything about Alice's and Bob's key.


## A perfect attack!

- What eve really wants is unitary that does this:
- $|0\rangle|0\rangle \rightarrow|0\rangle|0\rangle$
- $|1\rangle|0\rangle \rightarrow|1\rangle|1\rangle$
- $|+\rangle|0\rangle \rightarrow|+\rangle|+\rangle$
- $|-\rangle|0\rangle \rightarrow|-\rangle|-\rangle$
- If she had such a machine, she could learn everything about Alice's and Bob's key. Just wait until Alice announces the basis she used, and then measure in that basis!


## Clicker Question

- Suppose there was a unitary that did this:
- $|0\rangle|0\rangle \rightarrow|0\rangle|0\rangle$
- $|1\rangle|0\rangle \rightarrow|1\rangle|1\rangle$
- $|+\rangle|0\rangle \rightarrow|+\rangle|+\rangle$
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It would
A) Preserve inner products
B) Violate state normalization
C) Not be a unitary
D) Violate causality

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## Information gain/disturbance

## tradeoff

- This "perfect attack" is unphysical---it can't be implemented by unitary + measurements
- This is basically a consequence of the no-cloning theorem---you can't make copies of an unknown quantum state
- There's a more general phenomenon: if you do an operation that extracts even a little information about the state, it must disturb the state.
- In BB84, any such disturbance can be detected by Alice and Bob.
- Fancier analysis lets Alice and Bob figure out how much information about the key is leaked based on how much noise they observe.
- If it's not too much leaked, they can fix it up to be totally secure (hashing/privacy amplification)
- This fixing up is important because real systems will have noise, and would still like to be able to generate key.
- For us for now, just know that alice and bob can detect any attempt at eavesdropping, and can't be fooled into thinking they have secure key when they don't.


## BB84 aparatus



## QKD in space



## What we've learned

- A one-time pad lets Alice and Bob communicate securely even if Eve listens in.
- They cannot do this remotely with classical signals
- They CAN do this remotely with quantum signals.


## Next Class

- More Alice and Bob
- Dense Coding
- Please read 6.4
- Note: we are skipping 6.3 for now.


## Problem Set

- New problem set will be on canvas this afternoon.
- It may not be on the course website till tomorrow.


## Dense Coding

