Turing Machines, contd.

Lecture 10
Turing Machine

Finite alphabet

Read

Write

Move +1 or -1

Halt condition

Internal state (finite number)
TM for Decision Problems

\[ M = (Q, \Sigma, \Gamma, B, \delta, q_{\text{start}}, q_{\text{accept}}, q_{\text{reject}}) : \]

\( \Gamma \) is a finite tape alphabet.

- \( B \) or \( \square \) is the blank symbol (special symbol)
- \( \Sigma \) is a finite input alphabet \( \Sigma \subseteq \Gamma \setminus B \)

\( Q \) is a finite set of states

\[ q_{\text{start}} \in Q \text{ is the initial state} \]

\[ q_{\text{accept}}, q_{\text{reject}} \in Q \text{ accept/reject states} \]

Or maybe run forever

Transition function: \( \delta : Q \times \Gamma \text{ (read)} \rightarrow Q \times \Gamma \text{ (write)} \times \{ L, R \} \)
TMs: what we saw and will see

• They are quite tedious to program, but possible! (it’s the assembly language version)

• They can do anything a computer can do (copy, shift, add…)

• e.g. RAM
TMs: what we saw and will see

- Will see that a TM can simulate itself. Write a TM interpreter in TM!
- Universal TM.
Church-Turing Thesis:

“Any physically realizable model of computation is equivalent to a TM”

- More of a physical law than a math theorem.
- e.g. Python doesn’t have additional power over TM.
- sounds fancy but it says no more than “a Python interpreter can compute anything you can compute in Python”
TMs: what we saw and will see

- **Church-Turing Thesis:**
  
  “Any physically realizable model of computation is equivalent to a TM”

- There are models of computation not equivalent to TM, we won’t see them this semester.
Variants/Extensions

Adding more capabilities to TMs make them easier to program

But doesn’t change what TMs can do: whatever the new variant can do, can be simulated in the original variant (with a lot more steps, sometimes)
Extension: multiple tracks

$M$ can address any particular track in the cell it is scanning

Can simulate multiple tracks with a single track machine, using extra “stacked” characters:

- Single new character
**Extension: multiple tracks**

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4 tracks

\[
M: \delta(q, -,0,-,-) = (p, -, -, -, 1, R)
\]

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infinite tape →

"If in state $q$ reading 0 on second track, then go to state $p$, write 1 on fourth track, and move right"

Then in $M'$ \(\delta(q, \begin{array}{c} x \\ 0 \\ y \\ z \end{array}) = (p, \begin{array}{c} x \\ 0 \\ y \end{array}, R)\) for every $x, y, z \in \Gamma$
Extension: multiple tracks

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\[
M: \delta(q, -,0,-,-) = (p, -, -, -, 1, R)
\]

“If in state \( q \) reading 0 on second track, then go to state \( p \), write 1 on fourth track, and move right”

Transition function:

\[
\delta : Q \times \Gamma_1 \times \Gamma_2 \times \Gamma_3 \times \Gamma_4 \rightarrow Q \times \Gamma_1 \times \Gamma_2 \times \Gamma_3 \times \Gamma_4 \times \{ L, R \}
\]
Extension: multiple tracks

4 tracks

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$M$: $\delta(q, -,0,-,-) = (p, -, -, -, 1, R)$

“If in state $q$ reading 0 on second track, then go to state $p$, write 1 on fourth track, and move right”

Transition function:

$\delta : Q \times (\Gamma_1 \times \Gamma_2 \times \Gamma_3 \times \Gamma_4) \rightarrow Q \times (\Gamma_1 \times \Gamma_2 \times \Gamma_3 \times \Gamma_4) \times \{ L, R \}$
Extension: multiple tracks

Sometimes intuitively better with multiple tracks

e.g. assume I want to copy this string.
Extension: multiple tracks

Sometimes intuitively better with multiple tracks

 e.g assume I want to copy this string.

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Extension: multiple heads

Transition function:

\[ \delta : Q \times \Gamma^2 \rightarrow Q \times \Gamma^2 \times \{L, R\}^2 \]
Snapshot of simulation (2 heads)

Single move:
\[ \delta(q_1,1,0) = (q_2,0,0,R,L) \]
Snapshot of simulation (2 heads)

Single move:
δ(q₁,1,0) = (q₂,0,0,R,L)
Single move:
\[ \delta(q_1,1,0) = (q_2,0,0,R,L) \]

- Simulate with multiple tracks. Special mark on track 1 and 2 for head positions. Track 0 has input.
- Make sweeps over the entire tape
Single move: 
\[ \delta(q_1,1,0) = (q_2,0,0,R,L) \]

1) Scan to the right to find the mark on track 1, read the corresponding symbol from track 0 into our internal state, and then return to the left end of the tape.
Snapshot of simulation (2 heads)

Single move: 
\( \delta(q_1, 1, 0) = (q_2, 0, 0, R, L) \)

2) Using \( M \)'s transition function, the internal state records \( M \)'s next state, the symbol to be written by each head, and the direction to move each head.
Snapshot of simulation (2 heads)

Single move:
\[ \delta(q_1,1,0) = (q_2,0,0,R,L) \]

3) Scan to the right to find the mark on track \( i \), write the correct symbol onto on track 0, move the mark on track \( i \) one step left or right, and then return to the left end of the tape.
Snapshot of simulation (2 heads)

$M$

Single move:
$\delta(q_1,1,0) = (q_2,0,0,R,L)$

3) Scan to the right to find the mark on track $i$, write the correct symbol onto on track 0, move the mark on track $i$ one step left or right, and then return to the left end of the tape.
Snapshot of simulation (2 heads)

Single move:
\[ \delta(q_1,1,0) = (q_2,0,0,R,L) \]

• Subroutine!
• However, seriously slows down the process but we don’t care about running time right now
Extension: multiple tapes

$k$-tape TM

$k$ different (2-way infinite) tapes

$k$ different independently controllable heads

input initially on tape 1; tapes 2, 3, ..., $k$, blank.

Single move:

read symbols under all heads

print (possibly different) symbols under heads

move all heads (possibly in different directions)

go to new state
Extension: multiple tapes
Extension: multiple tapes
Can’t compute more with k tapes

Theorem: If $L$ is accepted by a $k$-tape TM $M$, then $L$ is accepted by some 1-tape TM $M'$. 

Idea: $M'$ uses $k$ tracks to simulate tapes of $M$.

$M' \text{ will use } 2k \text{ tracks to simulate tapes+heads of } M$
Convention for TM

Input tape (read only, finite)

Work tape (read/write)

Output tape (write only)
Convention for TM

More convenient!

- Output doesn’t clash with input
- Don’t have to clean work tape
- Just remember to copy what I need to output tape
Extension: 2-Way Infinite Tape

How to do it with one infinite direction?
2-Way Infinite Tape: Folding

Simulating it in the original TM variant:

Modify transitions:
Remember in control if +ve or -ve side of tape
(contents of 0 cell will be marked).

If positive: R → RR & L → LL
If negative: R → LL & L → RR
At 0: R → R & L → RR
At 1?
2-Way Infinite Tape: multiple tracks

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| -1 | -2 | -3 | -4 | -5 |   |   |   |   |   |   |   |   |   |   |
2-Way Infinite Tape: shifting
2-Way Infinite Tape: shifting

When the machine reads ⏯ write a blank, move right, write a ⏯, move right and proceed as if we had read a blank.

When the machine reads ⏯ shift the entire contents of the tape to the right. Move back to the ⏯, move right, write a blank and proceed as if we had read a blank.
2-Way Infinite Tape: shifting

When the machine reads ▶ shift the entire contents of the tape to the right. Move back to the ▶, move right, write a blank and proceed as if we had read a blank.
Subroutine calls

Mechanism for $M_1$ to “call” $M_2$ on an argument

Goal:

I need to be able to do two things:

- push(q): push the state in some stack, save it.

- pop(q): pop whatever state is on top of stack and make it current state.
Subroutine calls

Implement the Stack with a new tape

- For push, write a new symbol to stack and move R
- For pop read symbol, write blank, move head L
Subroutine calls

- Recursion (e.g. Fibonacci)
- Can take existing TMs and call them as subroutines.
- Call = jump to start state of the TM subroutine
- Halt = return
Random Access Memory (RAM)

- By definition can only access memory directly under the head.
- How to do associative memory?
- Memory is made up from pairs [key,value]
- $key \in \{0,1\}^*$, $value \in \{0,1\}^*$
Random Access Memory (RAM)

- Would like a subroutine that starts with “key” written at the beginning of a tape and ends with “value” written at the beginning of the same tape for any key a most one value
Random Access Memory (RAM)

\[
\begin{array}{c}
[\text{key}, \text{value}] \quad [\text{key}, \text{value}]
\end{array}
\]

Ram tape \( \Sigma = \{[\text{}], 0, 1\} \)

\[
\begin{array}{c}
\text{key}
\end{array}
\]

Address tape

\[
\begin{array}{c}
\text{value}
\end{array}
\]

Value tape
Random Access Memory (RAM)

- If I have an RAM also that runs in time $T(n)$, I can simulate it in one tape, one head, one track TM in time $T(n)^2$
Universal Turing Machine

- "Turing machine interpreter written in Turing machine".

- Just as the input to a Python interpreter is a string of Python source code, the input to our universal Turing machine U is a string $\langle M, w \rangle$ that encodes an arbitrary Turing machine M and a string w in the input alphabet of M.

- Given these encodings, U simulates the execution of M on input w; in particular,
  - U accepts $\langle M, w \rangle$ if and only if M accepts w.
  - U rejects $\langle M, w \rangle$ if and only if M rejects w.
Universal Turing Machine

- How to encode a Turing Machine as a binary string:
  - $01|\Gamma|01|\Sigma|01|Q|0[....]$ where $[....]$ is some encoding (brute force) of all possible transitions as pattern of bits.

- Encode the tape as a bit string: (e.g. tape alphabet of 3 symbols $\{a,b,c\}$)

```
0  1  0  1  0  1
```

:tape was bac